

Neural Information Processing Systems Foundation

Online Learning

For each round $t = 1, \ldots, T$, we pick a point $x_t \in K$, and an adversary picks a cost function f_t . We incur the loss $f_t(x_t)$. The regret of our strategy is the difference between our total loss and the total loss of the best fixed point in hindsight:

$$R(T) = \sum_{t=1}^{T} f_t(x_t) - \arg\min_{x \in K} \sum_{t=1}^{T} f_t(x).$$

The goal is to minimize the regret R(T).

Delays in Learning

The standard models assume that the adversary gives us the loss function f_t before we select the next point x_{t+1} . What if the feedback is delayed? For example:

- Online advertising algorithms serve many ads before finding out which ones are clicked.
- Online algorithms planning resource allocation in the cloud cannot wait for one batch job to end before launching the next.
- In finance, online learning algorithms managing portfolios are subject to information and transaction delays from the market.
- Optimization algorithms implemented in distributed or parallel environments suffer communication delays between asynchronous processors.

Selected References

- A. Kalai and S. Vempala. Efficient algorithms for online decision problems. J. Comput. Sys. Sci., 71:291–307, 2005. Extended abstract in Proc. 16th Ann. Conf. Comp. Learning Theory *(COLT)*, **2003**.
- M. Zinkevich. Online convex programming and generalized infinitesimal gradient ascent. In Proc. 20th Int. Conf. Mach. Learning (ICML), pages 928–936, 2003.

where D is the sum of delays. In both cases the algorithms are left essentially unmodified.

Online Learning with Adversarial Delays

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Abstract

We study the performance on standard online learning algorithm when the feedback is **delayed by an adversary**. We obtain:

- $O(\sqrt{D})$ regret bounds for online-gradient-descent, and
- $O(\sqrt{D})$ regret bounds for follow-the-perturbed-leader,

Convex setting:

Convex domain K, convex loss functions $\{f_t : \mathbb{R}^n \to \mathbb{R}\}$

Undelayed algorithm and regret bound:

$$x_{t+1} = \pi_K \left(x_t - \Theta\left(\frac{1}{\sqrt{T}}\right) f'_t(x_t) \right)$$

where π_K projects to nearest point in K.

$$\Rightarrow \sum_{t=1}^{T} f_t(x_t) \le \underset{x \in K}{\operatorname{arg\,min}} \sum_{t=1}^{T} f_t(x) + O(\sqrt{T}).$$

Delayed algorithm:

$$x_{t+1} = \pi_K \left(x_t - \Theta\left(\frac{1}{\sqrt{D}}\right) \sum_{s \in \mathcal{F}_t} f'_s(x_s) \right)$$

(!) Same as undelayed algorithm when $\mathcal{F}_t = \{t\}$

Delayed regret bound:

$$\sum_{t=1}^{T} f_t(x_t) \le \underset{x \in K}{\operatorname{arg\,min}} \sum_{t=1}^{T} f_t(x) + O(\sqrt{D})$$
(!) Matches undelayed regret bound

 \triangleright $O(\sqrt{D})$ regret bound for follow-the-lazy-leader, a variation of follow-the-perturbed-leader for switching costs.

Delayed Feedback Model

- can be used in round $t + d_t$.

[Zinkevich, 2003]



Undelayed algorithm and regret bound:

$$x_{t+1} = \underset{x \in K}{\operatorname{arg\,min}} c_0 \cdot x + \sum_{s=1}^t c_s \cdot x,$$

where $c_0 \sim [0, \Theta(\sqrt{T})]^n$ uniformly at random.

$$\Rightarrow \sum_{t=1}^{T} c_t \cdot x_t \le \underset{x \in K}{\operatorname{arg\,min}} \sum_{t=1}^{T} c_t \cdot x + O(\sqrt{T})$$

Delayed algorithm:

 $x_{t+1} = \arg\min c_0 \cdot x$ $x \in K$

where $c_0 \sim [0, \Theta(\sqrt{D})]^n$ uniformly at random. (!) Same as undelayed algorithm when $\mathcal{F}_t = \{t\}$

Delayed regret bound:

when D = T

$$\sum_{t=1}^{T} c_t \cdot x_t \leq \underset{x \in K}{\operatorname{arg\,min}} \sum_{t=1}^{T} c_t \cdot x + O(\sqrt{D})$$
(!) Matches undelayed regret bound when $D = T$

Extensions

 \triangleright $O(\sqrt{D})$ regret bound for online-mirror-descent, a generalization of online-gradient-descent and randomized expert selection by exponential weights.



▶ $d_t \in \mathbb{Z}^+$ denotes a non-negative **delay**. Feedback from round t is delivered at the end of round $t + d_1 - 1$ and

▶ $\mathcal{F}_t = \{u \in [T] : u + d_u - 1 = t\}$ denotes the set of rounds whose feedback appears at the end of round t.

• $D = \sum_{t=1}^{T} d_t$ denotes the sum of all delays. In the standard setting with no delays, D = T.

follow-the-perturbed-leader

Discrete domain K, cost vectors $\{c_t \in \mathbb{R}^n\}$

[Kalai and Vempala, 2005]

$$+\sum_{s=1}^{t}\sum_{r\in\mathcal{F}_s}c_r\cdot x$$